

# UIUC Mock Putnam Exam 3/2003

## Advanced Version

### Solutions

**Problem 1.** A *binary partition* of an integer  $n$  is a partition of  $n$  into parts of the form  $2^i$ ,  $i = 0, 1, \dots$  (with repetition allowed). For example,  $n = 5$  has the following binary partitions:  $5 = 4 + 1 = 2 + 2 + 1 = 2 + 1 + 1 + 1 = 1 + 1 + 1 + 1 + 1$ . Let  $b(n)$  denote the number of binary partitions of  $n$ ; the first few values of  $b(n)$  are  $b(1) = 1$ ,  $b(2) = 2$ ,  $b(3) = 2$ ,  $b(4) = 4$ ,  $b(5) = 4$ . Show that  $b(n)$  is even for all  $n \geq 2$ .

**Solution.** We first establish a recurrence for  $b(n)$ . Let  $b_k(n)$  denote the number of binary partitions with exactly  $k$  1's. Then  $b(n) = \sum_{k=0}^n b_k(n)$ . Furthermore,  $b_n(n) = 1$ , and if  $k < n$ , the number of binary partitions of  $n$  with exactly  $k$  1's is equal to the number of binary partitions of  $n - k$  containing no 1's, i.e., the number of partitions of  $n - k$  into parts of the form  $2, 2^2, \dots$ . But the latter number is only nonzero if  $n - k$  is even, and dividing through by 2 we see that in this case it is equal to the number of unrestricted binary partitions of  $(n - k)/2$ . Thus, for all  $k = 0, \dots, n - 1$  we have  $b_k(n) = b((n - k)/2)$  where  $b(x)$  is to be interpreted as 0 when  $x$  is not a positive integer. It follows that for  $n \geq 4$

$$\begin{aligned} b(n) &= \sum_{k=0}^{n-1} b_k(n) + b_n(n) = \sum_{k=0}^{n-1} b((n - k)/2) + 1 \\ &= \sum_{m=1}^{\lfloor n/2 \rfloor} b(m) + 1 = 2 + \sum_{m=2}^{\lfloor n/2 \rfloor} b(m). \end{aligned}$$

This shows that if the terms  $b(m)$ ,  $m = 2, 3, \dots, \lfloor n/2 \rfloor$ , are all even, then so is  $b(n)$ . Since  $b(2) = b(3) = 2$ , induction gives that  $b(n)$  is even for all  $n \geq 2$ .

**Problem 2.** Let  $P(x)$  be a polynomial of degree  $n$  satisfying  $P(k) = 1/k$  for  $k = 1, 2, \dots, n + 1$ . What is  $P(n + 2)$ ?

**Solution.** Let  $Q(x) = xP(x) - 1$ . Then  $Q(x)$  is a polynomial of degree  $n + 1$  with zeros at  $k = 1, 2, \dots, n + 1$ . By the factor theorem,  $Q(x)$  must be of the form (\*)  $Q(x) = c \prod_{k=1}^{n+1} (x - k)$ , where  $c$  is a constant (depending on  $n$ ). Also,  $Q(0) = 0 \cdot P(0) - 1 = -1$ , and setting  $x = 0$  in (\*) gives  $-1 = c \prod_{k=1}^{n+1} (-k) = c(-1)^{n+1}(n+1)!$ , so  $c = (-1)^n/(n+1)!$ . Hence  $Q(n+2) = c(n+1)! = (-1)^n$ , so  $P(n+2) = (1 + Q(n+2))/(n+2) = (1 + (-1)^n)/(n+2)$ , i.e.,  $P(n+2) = 0$  if  $n$  is odd, and  $P(n+2) = 2/(n+2)$  if  $n$  is even.

**Problem 3.** [1995 UIUC Undergrad Math Contest] Let  $f(n)$  be a non-negative function, defined on the set of nonnegative integers and satisfying  $f(n+m) \leq f(n) + f(m)$  for all  $n, m$ . Show that  $\lim_{n \rightarrow \infty} f(n)/n$  exists.

**Solution.** First, taking  $n = m = 0$  in the given inequality, we obtain  $f(0) \leq 2f(0)$ , which forces  $f(0) = 0$ . Next, let  $m$  be a positive integer. From the given inequality we obtain  $f(km) = f(m+(k-1)m) \leq f(m) + f((k-1)m)$  for all positive integers  $k$ . By induction, it follows that

$$f(km) \leq kf(m) \quad \text{for all } k \quad (*)$$

Now, set  $g(n) = f(n)/n$  for  $n \geq 1$ . Then (\*) implies  $g(n) \leq g(m)$  for all positive integers  $n$  that are multiples of  $m$ . On the other hand, if  $n = km + r$  with  $0 < r < m$ , then, using the given inequality we have

$$g(n) = \frac{f(n)}{n} \leq \frac{f(km) + f(r)}{n} = \frac{kmg(km) + rg(r)}{n} < g(m) + \frac{C_m}{n},$$

where  $C_m = \max_{r=1}^{m-1} rg(r)$  is a constant depending only on  $m$ . It follows that

$$\limsup_{n \rightarrow \infty} g(n) \leq g(m).$$

Since  $m$  was arbitrary, and the left-hand side does not involve  $m$ , we have

$$\limsup_{n \rightarrow \infty} g(n) \leq \inf_m g(m) \leq \liminf_{m \rightarrow \infty} g(m).$$

But this implies that the limsup and liminf of  $g(n)$  are equal, i.e., that  $\lim_{n \rightarrow \infty} g(n)$  exists, which is what we had to prove.

**Problem 4.** [1995 UIUC Undergrad Math Contest] Let  $c$  be a positive constant, let  $0 < x_1 < x_0 < 1$ , and for  $n \geq 1$  let  $x_{n+1} = cx_n x_{n-1}$ . Prove that there exists a positive real number  $\alpha$  such that the limit  $L = \lim_{n \rightarrow \infty} x_{n+1}/x_n^\alpha$  exists and  $0 < L < \infty$ .

**Solution.** We transform the sequence into one that satisfies the Fibonacci recursion. To this end, let first  $y_n = cx_n$ . Then  $y_n$  satisfies the recursion

$y_{n+1} = y_n y_{n-1}$  for  $n \geq 1$ , along with the initial condition  $0 < y_1 < y_0 < 1$ . Setting  $a_n = -\ln y_n$ , we then have  $a_{n+1} = a_n + a_{n-1}$  for  $n \geq 1$  and  $a_1 > a_0 > 0$ . The latter recurrence is the same as that satisfied by the Fibonacci numbers. Its characteristic equation,  $x^2 - x - 1 = 0$ , has roots at  $x = \phi$  and  $x = \bar{\phi}$ , where  $\phi = (1 + \sqrt{5})/2$  and  $\bar{\phi} = (1 - \sqrt{5})/2$ . By the general theory of second order linear recurrences, any solution to this recurrence is of the form  $a_n = c_1 \phi^n + c_2 \bar{\phi}^n$  with coefficients  $c_1$  and  $c_2$ . Since  $|\bar{\phi}| < 1$ , it follows that, as  $n \rightarrow \infty$ ,  $a_n - c_1 \phi^n \rightarrow 0$  and therefore  $a_{n+1} - \phi a_n \rightarrow 0$ . Hence  $y_{n+1}/y_n^\phi \rightarrow 1$ , and substituting  $y_n = c x_n$ , we get  $x_{n+1}/x_n^\phi \rightarrow c^{\phi-1}$  as  $n \rightarrow \infty$ . Thus, the limit  $L$  sought in the problem exists and is equal to  $c^{\phi-1}$ .